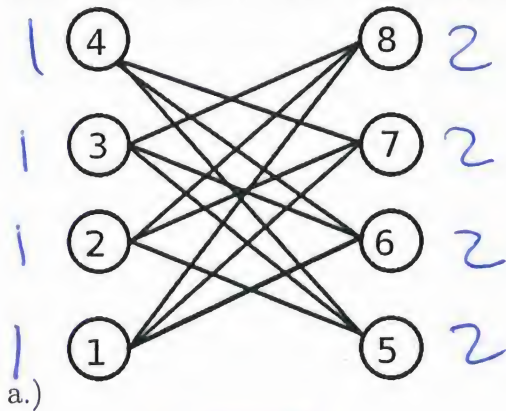


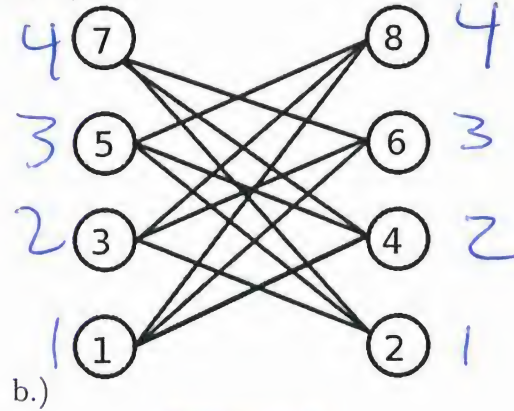
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Graph Theory Quiz 4 (14 June 2019)
Open book, open notes, open neighbor.

1. Using the greedy coloring algorithm, how many colors will result on the graph G below with the two given vertex orderings? (6 pts)



2



4

2. Tightly bound the possible chromatic numbers of G , $\chi(G)$. Justify your response using the bounds discussed in class.

G is bipartite

$\chi(G) = 2$

3. Is G color-critical? Justify your response.

No We've bounded $\chi(G) = 2$
In order to have $H = G - e$ s.t.
 $\chi(H) = 1$, G would need to
have only a single edge

4. Recall that C_n is color-critical for $n = \text{odd}$. Show that any graph G is k -color-critical for $\chi(G) = k = 3$ if and only if G is an odd cycle.

If G is odd cycle $\Rightarrow G$ is color-critical and $\chi(G) = 3$

- We demonstrated that all odd cycles have chromatic number of 3, in class with a greedy coloring argument.
 - Odd cycles are color-critical, as removing any edge creates a path, colorable w/ 2 colors
- G is color-critical and $\chi(G) = 3 \Rightarrow G$ is odd cycle
- Color-critical implies removing some e will decrease the chromatic number
 - G is not a cycle nor bipartite ($\chi(G) = 2$ as then)
- so G contains cycles, at least one of which is odd
- Consider removing a hypothetical e not in an odd cycle; as the odd cycle remains, the chromatic number is still 3
 - Hence, any such e must be in an odd cycle and there can only be one odd cycle
- $\Rightarrow G$ is an odd cycle \square